

# CS 294-92 Analysis of Boolean Functions

## Problem Set 2

Due: Feb 28, 2020, 7 PM

You are encouraged to discuss the problems and solve them in groups. However, the solutions are to be written up alone, listing all the collaborators.

1. **The Gravity Operator:** The “gravity operator”  $G_i$  in direction  $i$  maps Boolean functions  $f : \{-1, 1\}^n \rightarrow \{-1, 1\}$  to  $G_i f : \{-1, 1\}^n \rightarrow \{-1, 1\}$  as follows.

$$G_i f(x) = \begin{cases} \max\{f(x^{i \rightarrow -1}), f(x^{i \rightarrow +1})\}, & \text{if } x_i = +1 \\ \min\{f(x^{i \rightarrow -1}), f(x^{i \rightarrow +1})\}, & \text{if } x_i = -1 \end{cases}$$

Observe that by definition,  $G_i f$  is monotone in the  $i$ -th direction.

- (a) Show that if  $f$  is monotone in the  $j$ -th direction then so is  $G_i f$ , for any  $i, j \in [n]$ .
  - (b) Show that  $\Pr_{\mathbf{x} \sim \{-1, 1\}^n}[f(\mathbf{x}) = 1] = \Pr_{\mathbf{x} \sim \{-1, 1\}^n}[G_i f(\mathbf{x}) = 1]$ .
  - (c) Show that  $\mathbf{Inf}_i[f] = \mathbf{Inf}_i[G_i f]$  and  $\mathbf{Inf}_j[G_i f] \leq \mathbf{Inf}_j[f]$  for  $j \neq i$ .
  - (d) An edge-isoperimetric inequality on functions  $f : \{-1, 1\}^n \rightarrow \{-1, 1\}$  is an inequality of the form  $\mathbf{I}[f] \geq \phi(\Pr[f(\mathbf{x}) = 1])$  for some function  $\phi : [0, 1] \rightarrow \mathbb{R}$  (e.g., Poincaré’s Inequality correspond to  $\phi(\alpha) = 4\alpha(1 - \alpha)$ , and Harper’s Inequality correspond to  $\phi(\alpha) = 2\alpha \log_2(1/\alpha)$ ). Deduce that it suffices to prove edge-isoperimetric inequalities for monotone functions.
2. **Stability of Arrow’s Theorem & Extensions:** Recall that we proved in class Kalai’s Theorem stating that

$$\Pr_{\mathbf{x}, \mathbf{y}, \mathbf{z}}[(f(\mathbf{x}), f(\mathbf{y}), f(\mathbf{z})) \text{ is rational}] = \frac{3}{4} - \frac{3}{4} \mathbf{Stab}_{(-1/3)}[f]. \quad (1)$$

where for each  $i \in [n]$  independently  $(\mathbf{x}_i, \mathbf{y}_i, \mathbf{z}_i)$  are selected uniformly at random from all strings in  $\{-1, 1\}^3$  that are not-all-equal, and a result is called rational if the three bits are not-all equal.

- (a) Suppose  $1 - \varepsilon = \Pr_{\mathbf{x}, \mathbf{y}, \mathbf{z}}[(f(\mathbf{x}), f(\mathbf{y}), f(\mathbf{z})) \text{ is rational}]$ . Show that  $\mathbf{W}^1[f] \geq 1 - 9\varepsilon/2$ .
- (b) (Bonus) Extend the previous item to the case of three Boolean functions,  $f_1, f_2, f_3 : \{-1, 1\}^n \rightarrow \{-1, 1\}$  as long as  $\mathbf{E}[f_1] \mathbf{E}[f_2] + \mathbf{E}[f_2] \mathbf{E}[f_3] + \mathbf{E}[f_1] \mathbf{E}[f_3] \geq 0$ . Namely, show that if

$$1 - \varepsilon = \Pr_{\mathbf{x}, \mathbf{y}, \mathbf{z}}[(f_1(\mathbf{x}), f_2(\mathbf{y}), f_3(\mathbf{z})) \text{ is rational}]$$

then  $\mathbf{W}^1[f_1], \mathbf{W}^1[f_2], \mathbf{W}^1[f_3] \geq 1 - O(\varepsilon)$ .

- (c) Show that there exist functions  $f_1, f_2, f_3$  that are not dictators nor anti-dictators s.t.

$$\Pr_{\mathbf{x}, \mathbf{y}, \mathbf{z}}[(f_1(\mathbf{x}), f_2(\mathbf{y}), f_3(\mathbf{z})) \text{ is rational}] = 1.$$

### 3. Influences in Elections:

- (a) Ex. 2.9 in O’Donnell’s Book - “In 1965, the Nassau County (New York) Board used a weighted majority voting system to make its decisions, with the 6 towns getting differing weights based on their population. Specifically, the board used the voting rule  $f : \{-1, 1\}^6 \rightarrow \{-1, 1\}$  given by the weighted majority

$$f(x) = \text{sgn}(31x_1 + 31x_2 + 28x_3 + 21x_4 + 2x_5 + 2x_6).$$

Compute  $\mathbf{Inf}_i[f]$  for all  $i \in [6]$ .”

- (b) We define the composition of two Boolean functions  $f : \{-1, 1\}^n \rightarrow \{-1, 1\}$  and  $g : \{-1, 1\}^m \rightarrow \{-1, 1\}$ , denote by  $f \circ g : \{-1, 1\}^{mn} \rightarrow \{-1, 1\}$ , as

$$(f \circ g)(x) = f(g(x_1, \dots, x_m), g(x_{m+1}, \dots, x_{2m}), \dots, g(x_{m(n-1)+1}, \dots, x_{mn})).$$

What are the influences of the function  $\text{MAJ}_n \circ \text{MAJ}_n$ ? (asymptotically)

What are the influences of the function  $\underbrace{\text{MAJ}_3 \circ \text{MAJ}_3 \circ \dots \circ \text{MAJ}_3}_{k \text{ times}}$ ? (exactly)

- (c) (Bonus) Suppose you are a citizen in Randomland. Elections in Randomland are run between the two candidates: Positive and Negative. There are  $n$  citizens in Randomland. Out of them  $m$  citizens have decided to form the state of North-Random, and the others are “loners”. Assume  $m$  and  $n$  are odd. Elections are run as follows: everyone votes, but decide their vote uniformly at random. The majority winner in the state of North-Random gets  $m$  votes. Remaining votes by “loners” are counted directly. Then, based on the overall count the winner is decided. Estimate asymptotically in terms of  $m$  and  $n$  the influence of a loner versus the influence of a North-Random resident.

4. **Fourier Concentration of Decision Trees:** Let  $f : \{-1, 1\}^n \rightarrow \{-1, 1\}$  be computable by a decision tree of size  $s$  and let  $\varepsilon \in (0, 1]$ . Show that the spectrum of  $f$  is  $4\varepsilon$ -concentrated up to degree  $\log_2(s/\varepsilon)$ .

### 5. Lipschitz functions and Poincare’s Inequality.

- (a) We say that a function  $f : \{-1, 1\}^n \rightarrow \mathbb{R}$  is  $L$ -Lipschitz if for every neighboring points  $(x, y)$  on the hypercube  $\{-1, 1\}^n$ , we have  $|f(x) - f(y)| \leq L$ . Show that if  $f$  is  $L$ -Lipschitz then

$$\mathbf{E}_{\mathbf{x} \sim \{\pm 1\}^n} \mathbf{E}_{\mathbf{y} \sim \{\pm 1\}^n} [(f(\mathbf{x}) - f(\mathbf{y}))^2] \leq n \cdot L^2/2.$$

(Hint: Use Poincare’s Inequality).

- (b) **Extension to higher dimensions:** Suppose  $f : \{-1, 1\}^n \rightarrow \mathbb{R}^d$  is a function such that  $\|f(x) - f(y)\|_2 \leq L$  for all neighboring points  $(x, y)$  on the hypercube (where  $\|v\|_2 = \sqrt{\sum_{i=1}^d v_i^2}$  for  $v \in \mathbb{R}^d$ ). Show that

$$\mathbf{E}_{\mathbf{x} \sim \{\pm 1\}^n} \mathbf{E}_{\mathbf{y} \sim \{\pm 1\}^n} [\|f(\mathbf{x}) - f(\mathbf{y})\|_2^2] \leq n \cdot L^2/2.$$

- (c) Let  $DTSize : \{-1, 1\}^{2^n} \rightarrow \mathbb{N}$  be the function that on a given input  $x \in \{-1, 1\}^{2^n}$ , interprets  $x$  as a truth-table of a function  $f_x : \{-1, 1\}^n \rightarrow \{-1, 1\}$ , and outputs the size of the smallest decision tree computing  $f_x$ . Show that  $DTSize$  is  $n$ -Lipschitz.
- (d) Deduce that there exists a value  $s$  such that most (say 0.99 fraction) of the functions on  $n$  variables have decision tree size in the range  $[s, s + O(2^{n/2} \cdot n)]$ . (Note: You don’t need to determine  $s$ , just prove that it exists).